

Mine Blast Algorithm

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Abstract: Optimization is system by which we can do our work in most proficient manner as could be expected under the circumstances. This can be accomplished by lessening no. of counts specifically calculation or diminishing work done in any of the assignment. Mine blast calculation is another optimization calculation which is populace based calculation for taking care of obliged engineering issues. Computational exertion that is number of assessing functions are less and the aftereffects of the calculations are great. In excess of fifteen obliged issues are tackled by Mine blast calculation and results are progressively precise. A presentation and insight concerning the calculation just as one application which is utilizing mine blast for ideal outcomes is additionally depicted in detail. The name of the application is speed reducer issue in which the mine blast calculation is utilized to get the optimal aftereffects of the parameters which are optimal. This outcome is contrasted and different calculations and henceforth mine blast calculation demonstrated to be generally effective in giving ideal outcomes.

1. INTRODUCTION

Mine Blast Algorithm is an improvement calculation that is utilized to take care of many compelled engineering issues. Over the timeframe numerous calculations had been presented yet this calculation is demonstrated to give the better outcomes than other different algorithms. This is novel populace – based calculation. Numerous calculations came that uses gradient data for the assessment that leads to false data and a greater amount of the functionality evaluation. The

objective of optimization method is to decrease time, calculations, incorrectness, parameters on which calculations depend on. In this particular algorithm, the idea of mine blast is utilized.

Let's accept a circumstance where there is a mine field in which there are such a large number of mines plotted under the ground and to tidy up the mine field is the principle reason. One ideal point is taken in the beginning where the main mine is blasted and the shrapnel of that mine will blast different mines continuing toward this path we will get next ideal point. The principal ideal point from where the most extreme causalities caused is signified as X^* and the causalities caused are characterized by min and max work ($f(x)$) per X^* . Bombs of various sizes and powers are pointed under the ground. When a mine bomb explodes a significant number of its shrapnel pieces causes the causalities which are determined as the orientation and separation of every shrapnel is unique. One shrapnel can cause numerous causalities that shows the dangerous intensity of that bomb. These will prompt the clearance of the mine field. The solution area is field separated into the networks.

2. RELATED WORK

Mine Blast Algorithm provides the accuracy in results and leads to less function evaluations and the work is done in a specified domain. Algorithm has an initial point where the first mine is blasted known as first shot point. It is denoted as the A_0 . g is number of initial shot points (1,2,3,4...) and it is a user defined parameter. The population is generated by the first shot explosion which is denoted as the initial upper bound population (N_{pop}) and that is number of shrapnel (N_s). In this algorithm there are lower values which are provided in the problem. Now, the first shot point is calculated by the Equation (1):

$$A_0 = LB + \text{rand} * (UB - LB)$$

Where A_0 , LB and UB are first shot point, lower bound, upper bound and rand is uniformly distributed random number between 0 and 1

Now, the Equation (2):

$$A = \{A_m\}, y = 1, 2, 3, \dots, N_d$$

N_d is search space dimension which is equal to number of independent variables. Now, the

location of another mine bomb is given by Equation(3):

$$A_{(x+1)}^a = A_{(e(x+1))}^a + \exp(-\sqrt{((y_{(x+1)}^a)/(d_{(x+1)}^a)))} A_x^a, \quad x=0,1,2,3 \dots (3)$$

D is distance, m is distance. now the location of exploding bomb is given by Equation(4):

$$A_{(e(x+1))}^a = d_x^a \times \text{rand} \times \cos(\Theta), \quad x=0,1,2 \dots (4)$$

D is the distance to find next optimal point, theta is angle of the shrapnel pieces and is calculated as $360/N_s$. The distance and the direction of the shrapnel is calculated by Equation(5) and Equation(6):

$$d_{(x+1)}^a = \sqrt{((A_{(x+1)}^a - A_x^a)^2) + (F_{(x+1)}^a - F_x^a)^2} \quad x=0,1,2,3 \dots (5)$$

$$y_{(x+1)}^a = (F_{(x+1)}^a + F_x^a) / (A_{(x+1)}^a - A_x^a), \quad x=0,1,2,3 \dots (6)$$

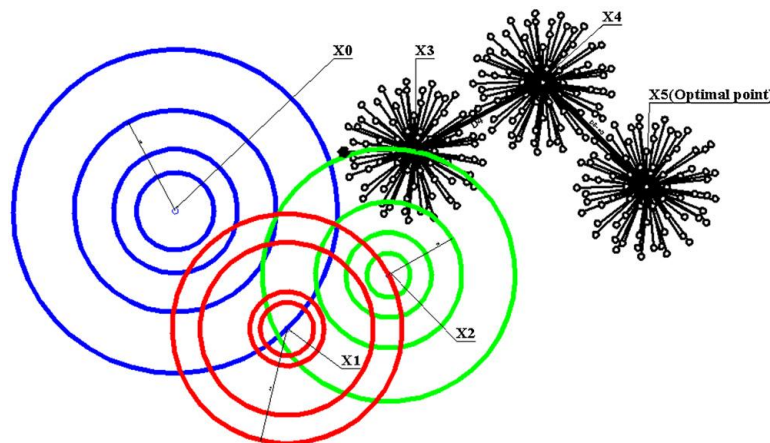


Fig.1 Exploration area of mines when the mines are blasted.

F is the function value for A. To check the exploration that is the range of the area or space smaller and larger distances the exploration factor(u) is used. The condition on the exploration factor is checked as:

If($u > k$)

Exploration(Equation(7) and(8))

Else

Exploration(Equation(4)(6) and(9)

End

K is iteration number index. for the solution space the Equation(7)and (8):

$$d_{(x+1)}^a = d_x^a \times (|\text{randn}|)^2, \quad x=0,1,2,\dots \quad (7)$$

$$A_{(x+1)}^a = d_{(x+1)}^a \times \cos(\Theta), \quad x=0,1,2 \dots \quad (8)$$

For the more accurate and rapid search we need to decrease the area is calculated by the Equation(9):

$$d_x^a = (d_{(x-1)}^a) / (\exp[\frac{f_0}{f_k}](k/\alpha)), \quad x=1,2,3,\dots \quad (9)$$

These were the parameters and the equations used in this algorithm. The functions of every equation is given below:

Equation(1): First shot point value.

Equation(2):current location of mine bomb

Equation(3):location of another mine

Equation(4):location of another mine to be used in previous equation

Equation(5):distance of shrapnel

Equation(6):direction of shrapnel pieces

Equation(7):new distance can be smaller or larger

Equation(8):new location of another mine bomb

Equation(9):reduced separation.

In the figure no.1 we can see that there are two procedures .In first procedure cycles are being done to get the ideal arrangement while in the second procedure the attention is determined to ideal focuses.

(i) Constraint Handling: There are a few guidelines dependent on which the outcomes are inferred, these are:

Rule1: over any infeasible arrangement, attainable arrangement is liked.

Rule2: on the off chance that there is little infringement in any infeasible arrangement, at that point it is considered as plausible arrangement. Rule3: in the event that there are two doable arrangements, at that point the arrangement having better target work esteem is taken. Rule 4: on

the off chance that there are two infeasible arrangements, at that point the arrangement with total of requirements is liked

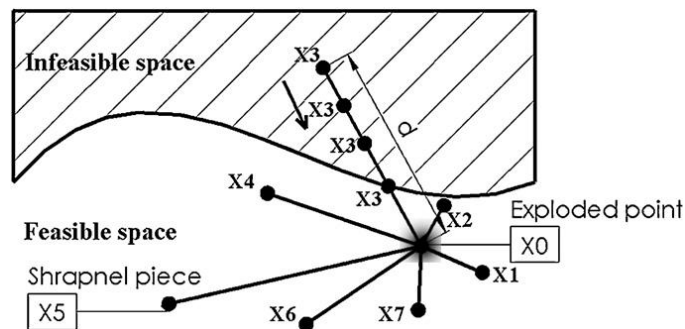


Fig 2: possible and infeasible region for limitation taking care of.

(ii) Termination criterion:

The calculation is ended when we get the best outcomes are determined. We can get the best outcomes in this calculation when the accompanying conditions are valid:

- a. Maximum number of emphasess.
- b. Less cpu time.

Steps of MBA:

Stage 1: we need to pick the underlying parameters N_s , u , an and most extreme number of cycles, where an is diminished removed.

Stage 2: the state of investigation that is ($u < k$) is checked.

Stage 3: presently in the event that the condition is fulfilled, at that point the condition no. (7) and (8) are applied to compute the separation of shrapnel. generally go to stage 10.

Stage 4: the bearing of the shrapnel is determined by Equation no (6).

Stage 5: presently, the shrapnel pieces are to be created and their improved area is likewise determined utilizing Equation no.(3).

Stage 6: for the produced shrapnel pieces the requirements are checked.

Stage 7: Temporal arrangement is spared which best shrapnel piece.

Stage 8: Then condition is watched that is if shrapnel piece have lower work esteem the best fleeting arrangement?

Stage 9: Now on the off chance that the condition is valid, at that point the best transient arrangement is traded with position of shrapnel piece.

Stage 10: Now the area and separation of the shrapnel pieces is determined by utilizing condition (4), (5) and come back to stage 4.

Stage 11: The separation of the shrapnel pieces is decreased by utilizing condition (9).

Stage 12: In the last advance the end rule is checked, on the off chance that the end standard is fulfilled, at that point the calculation is halted generally come back to Step 2.

2. APPLICATION:

Mine blast calculation is utilized in different applications with respect to the optimization of bracket structure, width factor and so on this calculation is demonstrated to give better outcomes when contrasted with different improvements like TRGA,GA,EC and so on. Here we will examine the application named Golinski's speed reducer issue manages the optimization during the plan of a gearbox. The target of the issue is to limit the heaviness of the speed reducer (then again, its volume) while keeping up the greatest proficiency of rotational speed of the two shafts. The speed reducer comprises of a rigging and pinion and particular shafts. There are 7 plan factors and two subsystems for this issue, apparatuses and shafts. Stress, geometric and redirection are the different imperatives included. The heaviness of the speed reducer incorporates both the heaviness of the riggings just as the heaviness of the poles.

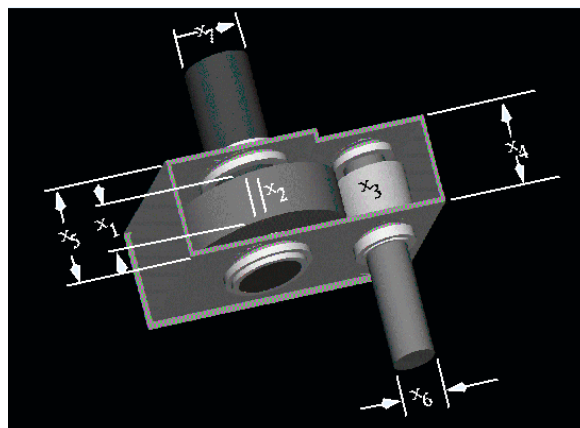


Fig 3. Round and hollow blade heat sink

4. PROBLEM DEFINATION

The main goal of Speed Reducer Golinski's issue is to find the base of a rigging box volume a (and, thus, its base weight), subject to a few limitations. There are seven plan factors, $q_1 - q_7$, which speak to

q_1 width of the rigging face, cm

q_2 teeth module, cm

q_3 number of pinion teeth

q_4 shaft 1 length between orientation, cm

q_5 shaft 2 length between heading, cm

q_6 breadth of shaft 1, cm

q_7 breadth of shaft 2, cm

5. RESULTS

The speed reducer issue is utilized in numerous different sorts of utilizations and is one of the standard benchmark issues in auxiliary advancement. There are numerous metaheuristic calculations that have been created to take care of this issue, these strategies can't ensure worldwide optimality of the arrangement. In any case, various break points are used in the linearization procedure for arriving at a worldwide approximate solution with a comparatively less errors, and much CPU time is required to settle the reformulated model. Consequently, this investigation likewise embraces the range reduction procedure to improve the computational productivity. Contrasted with various metaheuristic strategies, this investigation ensures the worldwide optimality of the solution.

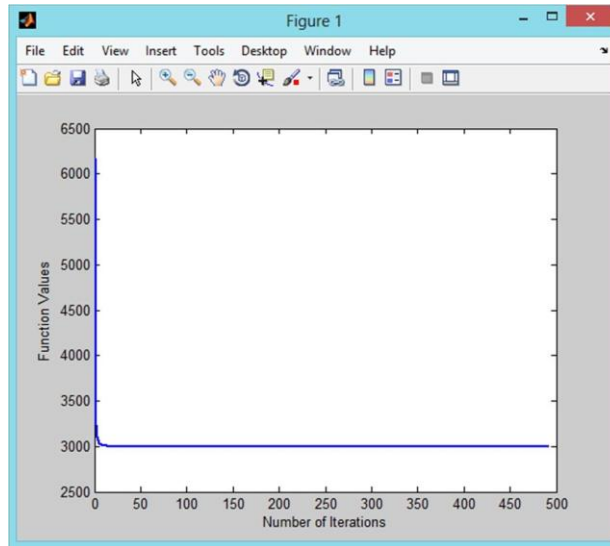


Table 2: shows the number of iterations and the corresponding function values.

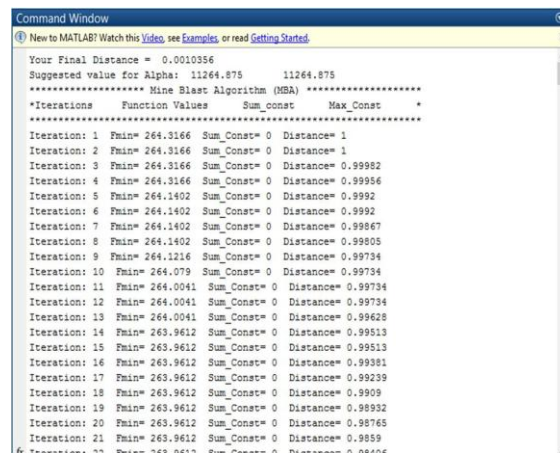
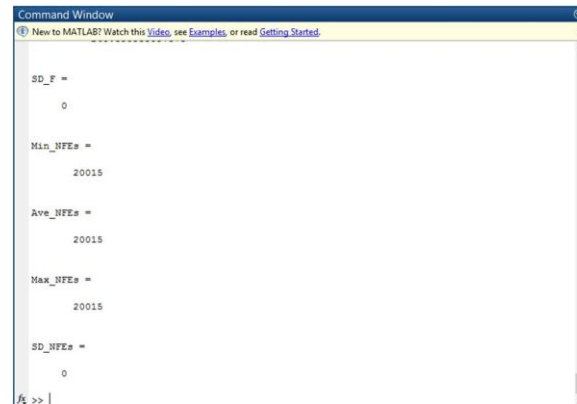


Fig 4: implementation of the mine blast algorithm in matlab.



Fig 5: Showing the values of the functions.**Fig:6: Showing the values of various functions of number of function evaluation.**

Hence the mine blast proved to give most efficient results in this application. As compared to other algorithms and most optimum results are given by the mine blast algorithm.

6.CONCLUSION

This paper presents the method and procedure of mine blast algorithm. The algorithm have very less number of function evaluation and the accuracy is more in the results. The basic concept of the algorithm is explosion of the mines in mine field and the main aim is to clear the mine field. Many iterations are carried out in this algorithm so as to get accurate and efficient results. The algorithm consists of 12 steps in which various equations which are easy to execute and more than fifteen constrained optimization and engineering problems are being solved and the results obtained by them gives better solutions than other optimizers terms of number of function evaluation(computational cost)and objective function for every problem. The quality of solution and computational efficiency obtained by Mine Blast Algorithm depends on the nature and complexity of underlined problem and is true for most meta heuristic methods.

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